111_015 Blockchain-2

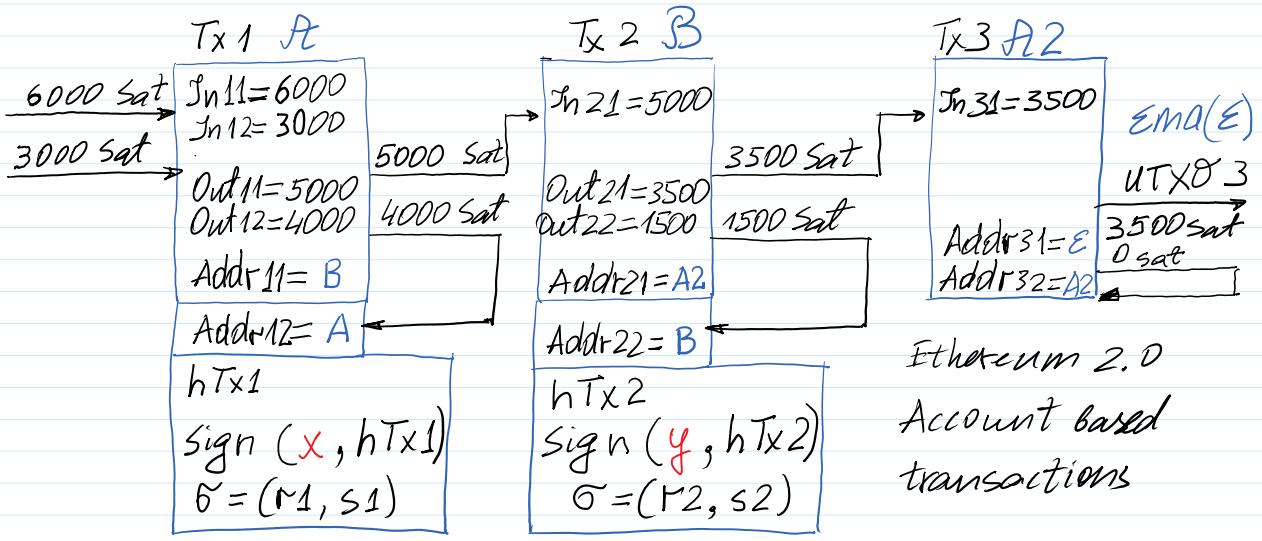
Mini https:
Merkle Tree. Paaiškinimas-diagrama-Google drive.

https://docs.google.com/spreadsheets/d/187fYuJTwltLJ4bG4tvQrNY8DXpR7fkRR/edit?usp=sharing&ouid=111502255533491874828&rtpof=true&sd=true

https://docs.google.com/document/d/11Bwk8HXLvjvzAEImcRiFcacwnrrz0lBs/edit?usp=sharing&ouid=111502255533491874828&rtpof=true&sd=true

2 Dalis. Reikia žodžiu atsakyti dėstytojui į vieną iš sekančių teorinių klausimų, remiantis paskaitų medžiaga:
2.1. Apklasis e-parašas ir jo panaudojimas.
2.2. Aklojo e-parašo demaskavimas ir jo panaudojimas.
2.3. Cut-and-Choose paradigma.
2.4. Atsitiktinė Identifikacijos Eilutė - Random Identification String (RIS).
2.5. Sukčiavimo nustatymas panaudojant RIS.
2.6. Pagrinddinės tranzakcijų sudėtinės dalys UTxO blokų grandinėje ir kriptografinių metodų panaudojimas .
2.7. Pagrindinės blokų sudėtinės dalys, kaip gaunama blokų grandinė, kuo paremtas jos saugumas ir blokų kasyba.
Atsakymas vertinamas maksimaliu balu 2.

Block structure - Unspent Transactio Output (UTxO) model



$Tx1 = '1 : In11 = 6000 || In12 = 3000 || Out11 = 5000 || Out12 = 4000 || Rec1 = B || Rec2 = A'$

$Tx2 = '2 : In21 = 5000 || Out21 = 3500 || Out22 = 1500 || Rec1 = A2 || Rec2 = B'$

$Tx3 = '3 : In31 = 3500 || Out31 = 3500 || Out32 = 0 || Rec1 = E || Rec2 = A2'$

Transaction template:
Tx_N = 'Tx_N:In11=...||In12=...||Out11=...||Out12=...||Rec1=...||Rec2=...'

Transactions:
Tx_1 = 'Tx_1:In11=6000||In12=3000||Out11=5000||Out12=4000||Rec1=B||Rec2=A'
Tx_2='Tx_2:In21=5000||Out21=3500||Out22=1500||Rec1=A2||Rec2=B'

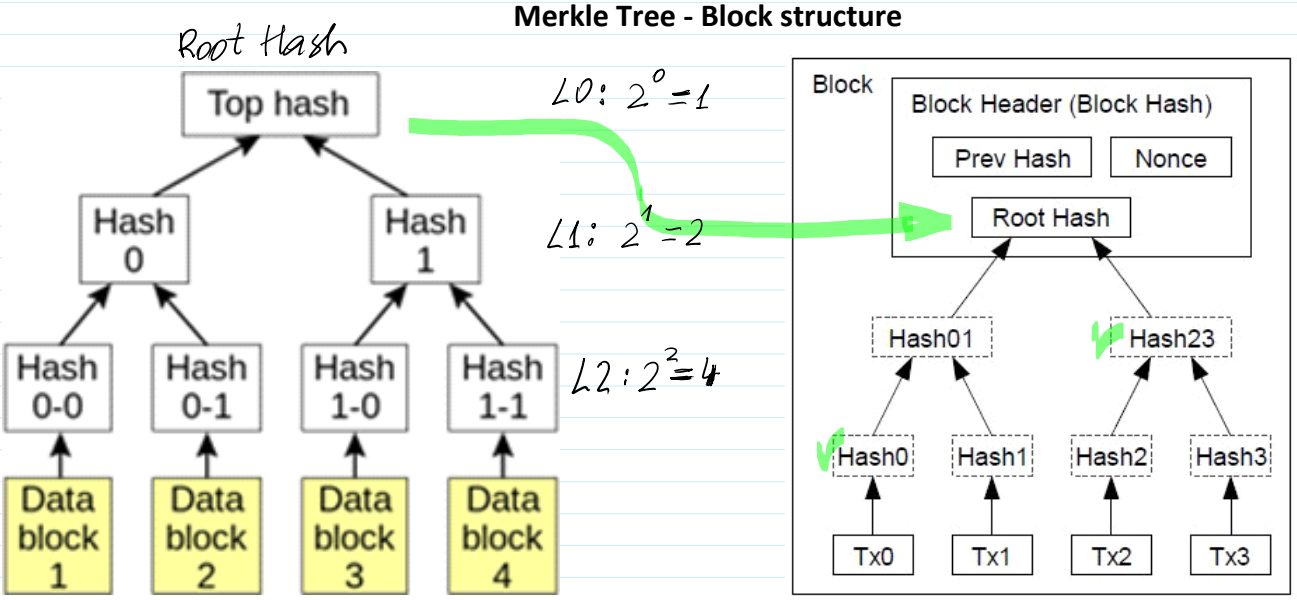
Tx_3='Tx_3:In31=3500||Out31=3500||Out32=0||Rec1=E||Rec2=A2'

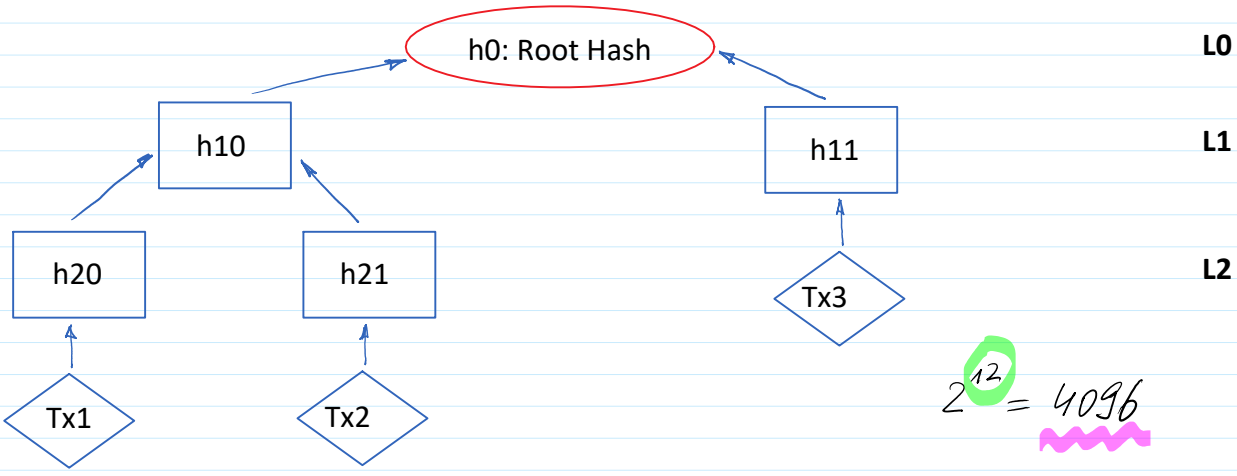
>> hTx_1=h28('Tx_1:In11=6000||In12=3000||Out11=5000||Out12=4000||Rec1=B||Rec2=A')
hTx_1 = 996BB7C
>> hTx_1=h28(Tx_1)
hTx_1 = 996BB7C

>> hTx_2=h28('Tx_2:In21=5000||Out21=3500||Out22=1500||Rec1=A2||Rec2=B')
>> hTx_2=h28(Tx_2)
hTx_2 = 977D75B

>> hTx_3=h28('Tx_3:In31=3500||Out31=3500||Out32=0||Rec1=E||Rec2=A2')
>> hTx_3=h28(Tx_3)
hTx_3 = 9201218

Merkle Tree - Block structure





>> h20=h28(hTx_1)

>> h20=h28(hTx_1)
h20 = 996BB7C
>> h21=h28(hTx_2)
h21 = 977D75B
>> h11=h28(hTx_3)
h11 = 9201218

>> h10=h28('996BB7C||977D75B')
h10 = 9201218

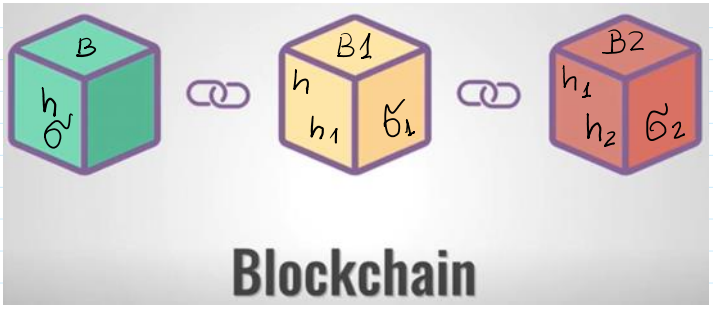
Root Hash: h0

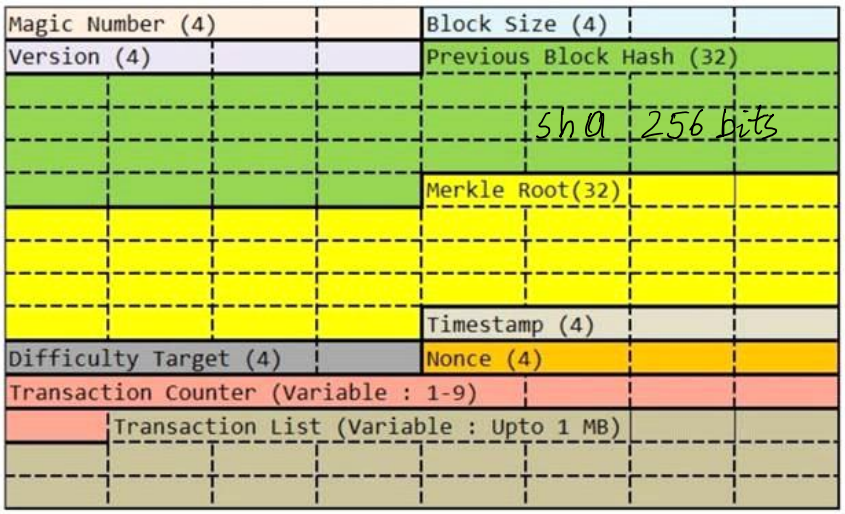
>> h0=h28('9201218||9201218')
h0 = **08E7C34**

Python : sha256

h20: 5B5412B
h21: D5C895A
h11: FEC59B7

h10: 625A41F
h0: **60BA3B5**





Block size = 4 Bytes
4 Bytes × 8 bits = 32 bits
Block have
$2^{32} - 1$ = 4294967295
In ASCII encoding
8 bits represents
1 symbol a, b, c, ... #
Block represents
536 870 912 symbols

Difficulty Target (DT): defines the complexity of block mining. In our simulation DT we will choose to find h-value of mining (mined block) having only 1 leading hexadecimal digit equal to 0_h.

h28('RootHash␣PrevHash␣73732763 1') =

>> sha256('RootHash PrevHash 737327631')
ans = F4AE534CD226FAF799 8C8424B348E020BA80639A687E93A0B8C5130ED C51E6DE

C51E6DE

>> sha256('RootHash PrevHash 737327632')
ans = B856211DF2EE15E30AB770C1A43CE014ECFE573182AFD885B28D96854DBC5F21
>> sha256('RootHash PrevHash 737327633')
ans = 9C18C764E347A58E57AC3F7A3C2874D5889A0E802699FEA47EEFF8C03BFEDA69

DT: to mine a block it is needed to find h-value having leading zero in hexadecimal format: C51E6DE 0XXXXXX

F 1111

$6 \times 4 = 24$ bits

h-value is computed su >> h28( ) → 7 hex numbers

What probability to mine a block?

The number of possible h-values of 28 bits: 2^{28} >> 2^28 ans = 268 435 456

The number of adequate h-values: 2^{24} >> int64(2^24) ans = 16777216

$$Pr\{\text{to Mine}\} = \frac{2^{24}}{2^{28}} = \frac{1}{2^4} = \frac{1}{16}$$

DT: two leading hex number = 00 00XXXXX

$5 \times 4 = 20$

The number of adequate h-values: 2^{20}

$$Pr\{\text{to Mine}\} = \frac{2^{20}}{2^{28}} = \frac{1}{2^8} = \frac{1}{256}$$

DT: two leading hex number = 000 000XXXX

$4 \times 4 = 16$

$$Pr\{\text{to Mine}\} = \frac{2^{16}}{2^{28}} = \frac{1}{2^{12}} = \frac{1}{4096}$$ >> 2^12 ans = 4096

$$Pr\{\text{to Mine}\} = \frac{1}{2^{28}} = \frac{1}{268\,435\,456}$$ >> 2^28 ans = 268 435 456

The probability to mine a block, e.g. in Bitcoin when DT: is to find SHA256 value having 18 leading zeroes

gas

1 Eth = 10^{18} gas

>> sha256('RootHash PrevHash 737327631')
ans = F4AE534CD226FAF799 8C8424B348E020BA80639A687E93A0B8C5130EDC51E6DE

000000000000000000 XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX

The number of possible h-values having 256 bits is 2^{256}.

The numbe of adequate h-values of SHA 256 is

$256 - 18 \cdot 4 = 256 - 72 = 184$ bits, that are represented 46 hex. num.

The number of adequate values is 2^{184}.

$$Prob\{to\ mine\} = \frac{2^{184}}{2^{256}} = 2^{184-256} = 2^{-72}$$

$1\,K = 2^{10} = 1024$

$1\,M = 2^{20} = \dots$

$1\,G = 2^{30} = \dots$

$1\,T = 2^{40} = \dots$

$2^{72} \sim 4\,GT = 4 \cdot 2^{30} \cdot 2^{40} = 2^2 \cdot 2^{30} \cdot 2^{40} = 2^{72}$.

4 722 366 482 869 645 213 696

Number of trials $N = 1T \cdot 1G \cdot 2^2 = 4 \cdot 2^{40} \cdot 2^{30}$.

Total net capacity $Cap \sim 2000$ Th/sek

$$Time\ T = \frac{N}{Cap} = \frac{4 \cdot 2^{40} \cdot 2^{30}}{2000 \cdot 2^{40}} \approx \frac{4 \cdot 2^{30}}{2^{11}} = 4 \cdot 2^{19}$$

```
>> T=int64(4*2^19)
T = 2097152
>> Tval=T/3600
Tval = 583
>> Tdien=Tval/24
Tdien = 24
```

Private blockchain ⟷ Public blockchain

Monero blockchain: Transactions sums → confidential → verifiable

Sender, Receiver → anonymous

How to realize confidential & verifiable transactions.